

3.2 Equivalence Relations

Def: Let R be a relation on a set A . * R is reflexive if $aRa \forall a \in A$.

* R is symmetric if whenever aRb , then bRa

* R is transitive if whenever aRb and bRc , then aRc .

Ex] \subseteq on \mathbb{R} * Is \subseteq reflexive? Yes; $a \subseteq a \forall a \in \mathbb{R}$

* Is \subseteq symmetric? No; $1 \subseteq 3$ but $3 \notin 1$.

* Is \subseteq transitive? Yes; If $a \subseteq b$ and $b \subseteq c$, then $a \subseteq c$.

Ex] \subseteq on $P(S)$ = set of all subsets of S

* reflexive? Yes; $A \subseteq A \forall A \in P(S)$

* symmetric? No: $S = \{1, 2, 3\}$, $A = \{1\}$, $B = \{1, 2\} \Rightarrow A \subseteq B$ but $B \notin A$

* transitive? Yes: If $A \subseteq B$ and $B \subseteq C$, then $A \subseteq C$.

Ex] = on \mathbb{Z}

Clearly reflexive, symmetric, and transitive.

Ex: $A = \{1, 2, 3\}$

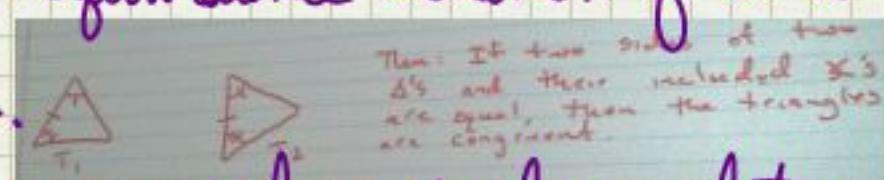
$R = \{(1, 1), (2, 2), (3, 3), (1, 2), (2, 3)\}$

* Reflexive? Yes

* Symmetric? No: $(1, 2) \in R$, but $(2, 1) \notin R$

* Transitive? No: $1 R 2$ and $2 R 3$, but $1 R 3$

Def: A relation R on a set A is an equivalence relation if it is symmetric, reflexive, and transitive.



Notation: We will use \sim to denote a general equivalence relation.

If $a \sim b$, we'll say a is equivalent to b . * symmetric? Yes: if $\Delta_1 \cong \Delta_2$, then $\Delta_2 \cong \Delta_1$.

Ex] Let $A = \{\text{triangles in the Euclidean Plane}\}$ * reflexive? Yes; given any $\Delta \in A$,

Consider geometric congruent, \cong . * transitive: yes: if $\Delta_1 \cong \Delta_2$ & $\Delta_2 \cong \Delta_3 \rightarrow \Delta_1 \cong \Delta_3$

Ex] Define \sim on \mathbb{Z} by $a \sim b \leftrightarrow 6|(b-a)$. Prove that \sim is an equivalent Rel.

Reflexivity: Let $a \in \mathbb{Z}$. Since $0=0 \cdot 6$, we know that $6|0$. But $a-a=0$, so $6|(a-a)$. Thus $a \sim a$.

Symmetric: Suppose $a \sim b$. Then $6|(b-a)$, so $\exists k \in \mathbb{Z}$ s.t. $b-a=6k$. Then $a-b=6(-k)$. This implies $6|(a-b)$, so $b \sim a$.

Transitivity: Suppose $a \sim b$ and $b \sim c$. Then this means $6|b-a$ and $6|c-b$.

Thus $\exists k, l \in \mathbb{Z}$ s.t. $b-a=6k$ and $c-b=6l$. Now $c-a=(c-b)+(b-a)=6k+6l=6(k+l)$, so $6|(c-a)$ and thus $a \sim c$.

Def: Let \sim be an equivalence relation on a set A and let $x \in A$.

The equivalence class of x , denoted \bar{x} is the set $\bar{x} = \{y \in A \mid x \sim y\}$

Ex] Recall \sim defined on \mathbb{Z} by $a \sim b$ iff $6|(b-a)$.

$$\text{so } \bar{0} = \{\dots, -12, -6, 0, 6, 12, \dots\}$$

$$1 \sim \bar{1} = \{\dots, -11, -5, 1, 7, 13, 19, \dots\}$$

similarly,

$$\bar{2} = \{\dots, -10, -4, 2, 8, 14, \dots\}$$

Note: Every element a is in \bar{a} . Either $\bar{a}=\bar{b}$ or $\bar{a} \cap \bar{b}=\emptyset$. $\bigcup_{a \in \mathbb{Z}} \bar{a} = \mathbb{Z}$

$$\bar{2} = \{\dots, -10, -4, 2, 8, 14, \dots\}$$

$$\bar{3} = \{\dots, -9, -3, 3, 9, 15, \dots\}$$

$$\bar{4} = \{\dots, -8, -2, 4, 10, 16, \dots\}$$

$$\bar{5} = \{\dots, -7, -1, 5, 11, 17, \dots\}$$

$$\bar{6} = \bar{0}$$

$$\bar{7} = \bar{1}$$

so $\bar{0}, \bar{1}, \bar{2}, \bar{3}, \bar{4}, \bar{5}$, and $\bar{6}$ are all of the equiv. classes of \sim .

Theorem: Let \sim be an equivalence relation on A , and let $x, y \in A$.

Then $x \in \bar{x}$.

$$\bigcup_{x \in A} \bar{x} = A$$

$\bar{x} = \bar{y}$ or $\bar{x} \cap \bar{y} = \emptyset$.

Def: Let A be a set and let $\{B_\alpha\}_{\alpha \in I}$ be a family of subsets of A satisfying

1. $B_\alpha \neq \emptyset \forall \alpha \in I$
2. $B_\alpha \cap B_\beta = \emptyset$ or $B_\alpha = B_\beta$
3. $\bigcup_{\alpha \in I} B_\alpha = A$

Then $\{B_\alpha\}_{\alpha \in I}$ is called a partition of A .

Ex] Using \sim as defined on \mathbb{Z} previously, $\{\bar{0}, \bar{1}, \bar{2}, \bar{3}, \bar{4}, \bar{5}\}$ is a partition of \mathbb{Z} .

Theorem: (Equivalent class theorem) 1. Let \sim be an equivalent rel. on a set $A \neq \emptyset$. Then $C = \{\bar{x}\}_{x \in A}$ forms a partition of A . 2. Define \sim' on A by $a \sim' b$ iff $a, b \in B_\alpha$ for some $\alpha \in I$. Then \sim' is an equivalent relation.

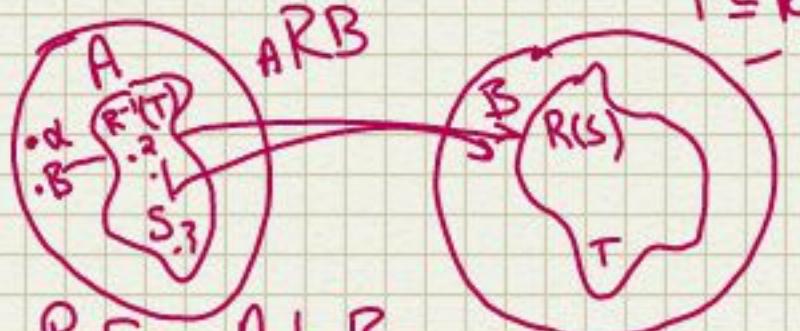
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Fact: Let R_1 and R_2 are equiv. relations on A .

1. $R_1 \cap R_2$ is an equiv. rel.

2. $R_1 \cup R_2$ is reflexive and symmetric
(but not necessarily transitive)

For R from A to B w/ $S \subseteq \text{Dom}(R)$
Prove $S \subseteq R^{-1}(R(S))$



R from A to B
 $S \subseteq A$ $T \subseteq B$

$S = \{1, 2, 3\}$

$$R = \{(1, 3), (1, 4), (2, 4), (3, 3), (1, 7, 0)\}$$

$$R(S) = \{3, 4\}$$

$$R^{-1}(\{3, 4\}) = \{1, 2, 3\}$$

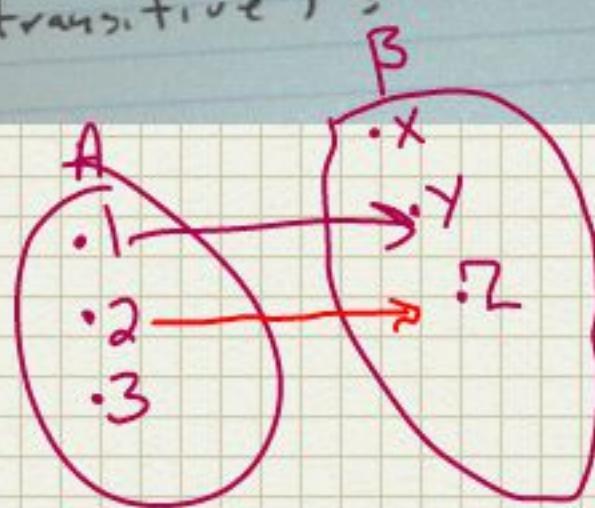
$$S = \emptyset \quad R(\emptyset) = \emptyset$$

$$R^{-1}(\emptyset) = \emptyset$$

$$S = \{18\} \quad R(S) = \emptyset$$

$$R^{-1}(\emptyset) = \emptyset$$

$$\emptyset \neq S \square$$

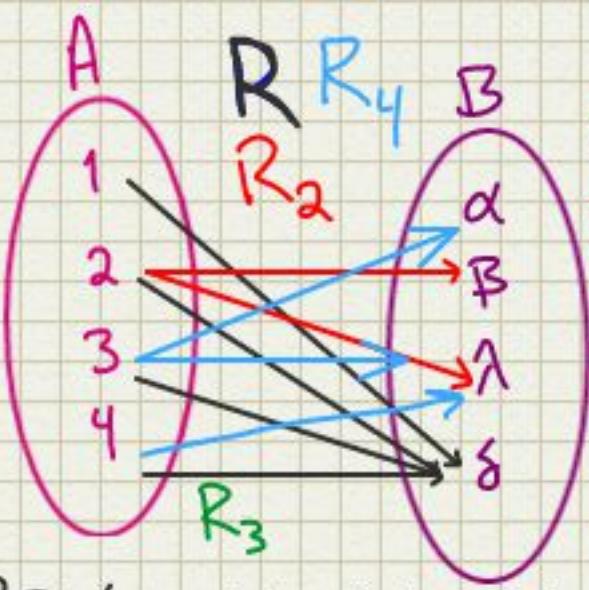


$$S = 2, 3 \quad R = (1, y)(2, z)$$

$$R = (1, y) \quad \text{Dom } R = \{1, 2\}$$

$$\text{Ran}(R) = \{y, z\}$$

$$R(S) = y, z$$



3a] Now $(x, y) \in R$ where $x \in S$
 $y \in R(S) \rightarrow x \in R^{-1}(R(S))$

$$R_3 = \{\}$$

$$\text{Dom } R_3 = \emptyset = S$$

$$\text{Ran } R_3 = \emptyset = T$$

$$R'^{-1}(R_3(S)) = R'^{-1}(\emptyset) = \emptyset = S$$

$$R_4 = \{(3, \alpha), (3, \gamma), (4, \gamma)\}$$

$$\text{Dom } R_4 = \{3, 4\} = S$$

$$\text{Ran } R_4 = \{\alpha, \gamma\} = T$$

$$R_4^{-1}(R_4(S)) = R_4^{-1}(\{\alpha, \gamma\}) = \{3, 4\} = S$$

$$ARB = \{(1, \delta), (2, \delta), (3, \delta), (4, \delta)\}$$

$$\text{Dom } R = \{1, 2, 3, 4\} = S$$

$$\text{Ran } R = \{\delta\} = T$$

$$R^{-1}(R(S)) = R^{-1}(\delta) = \{1, 2, 3, 4\} = S$$

$$ARB_2 = \{(2, \beta), (2, \lambda)\}$$

$$\text{Dom } R_2 = \{2\} = S; \text{Ran } R_2 = \{\beta, \lambda\} = T$$

$$R_2^{-1}(R_2(S)) = R_2^{-1}(\beta, \lambda) = \{2\} = S$$

#11)

Let $a \in A$. Since $\text{Dom}(R) = A$, then $a \in \text{Dom}(R)$. $(a, b) \in R$ for some $b \in B$.

Then $(b, a) \in R$. By symmetry, $(a, a) \in R$. For $\text{Dom}(R) = A$, (a, a)

$(a, b)R(c, d) \leftrightarrow (a-c) \in \mathbb{Z}$ and $(b-d) \in \mathbb{Z}$

$(c, d)R(c-d) \rightarrow c-c=0 \checkmark$ HW5 #3. Let $R = \{(a, b) \in \mathbb{Z} \times \mathbb{Z} \mid |a-b| < 5\}$

$$R = \{ \dots, (-99, -98), \dots, (-4, 0), (0, 4), (98, 99), (99, 98), \dots \}$$

$(a, b)R(c, d) \leftrightarrow (c, d)R(a, b)$

$$R = \{ \dots, (-3, 3), (99, 99), (0, 0), \dots \} \quad \begin{matrix} \text{symmetric} \\ (a, a) \in R \forall a \in \mathbb{Z} \end{matrix}$$

$a-c \in \mathbb{Z} \leftrightarrow c-a \in \mathbb{Z} \checkmark$

$$R = \{ (8, 4), (4, 0) \} \quad \begin{matrix} \text{reflexive} \\ |8-0| < 5 \end{matrix} \quad \begin{matrix} a-a=0 \\ \text{not transitive} \end{matrix}$$

$$R = \{ (3, 7), (7, 3) \} \wedge 7 \neq 3 \therefore \text{not antisymmetric}$$

3.3 Properties of relations on a set:

Ex \subseteq on \mathbb{Z}

Def: Let R be a relation on a set A . Irreflexive? No: $I \subseteq I$

• R is irreflexive if $aRa \forall a \in A$

Asymmetric? No: $I \subseteq I$ does not imply $I \neq I$.

• R is asymmetric if aRb implies bRa . Antisymmetric? Yes: $a \leq b \wedge b \leq a \Rightarrow a = b$

• R is antisymmetric if aRb and bRa implies $a = b$. Ex] \subset on \mathbb{Z}

Ex \subseteq on $P(S)$

Irreflexive? No: If $A \subseteq S$, then $A \subseteq A$

Irreflexive? Yes: $A \neq A \forall A \in P(S)$

Asymmetric? No: $A \subseteq A$ does not imply $A \neq A$

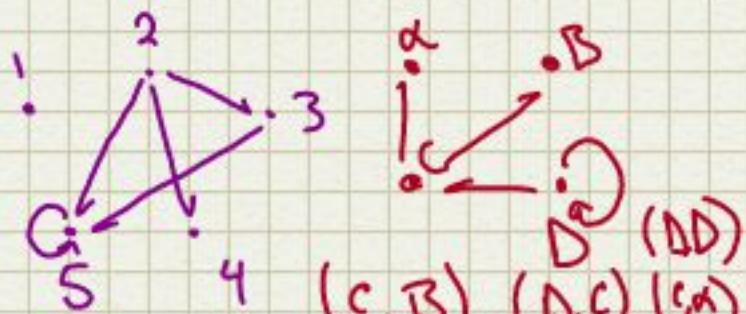
Asymmetric? Yes: If $a < b$, then $b \neq a$

Antisymmetric? Yes. If $B \subseteq A \wedge A \subseteq B$, $A = B$

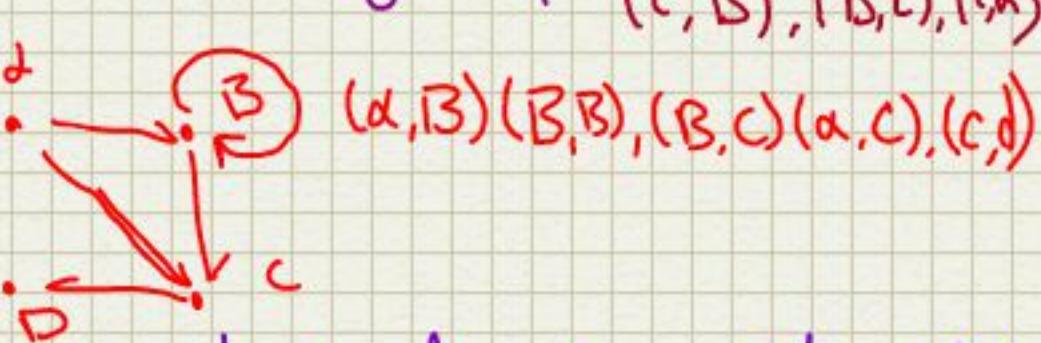
Let $A = \{1, 2, 3, 4, 5\}$
and $R = \{(2, 3), (2, 4), (3, 5), (2, 5), (5, 5)\}$

Ex \subset on $P(S)$

Irreflexive? Yes: If $A \subseteq S$, then $A \neq A$



Asymmetric? Yes: If $A \subset B$, then $B \not\subset A$



Digraphs of Relations

A digraph of a relation R on a set A is a way to visualize the properties of R .

Nodes of the digraph corresponds to the elements of A .

Ex: Make a digraph of $R = \{(a, a), (c, c), (a, b), (b, a), (b, d), (c, d), (b, c)\}$ on $A = \{a, b, c, d\}$.

Reflexive? X
double arrows
a \rightarrow b; b \rightarrow c
no loops
irreflexive? X
symmetric? X
a \rightarrow b; b \rightarrow c
transitive? X
no a \rightarrow c
asymmetric? X
antisymmetric? X
 $a \neq c$

3.4: Orderings

Def: A relation R on a set A is a partial order of A if R is reflexive, antisymmetric, and transitive.

Notation: \prec will denote general partial order.

Examples: \leq on \mathbb{Z} , $P(S)$, \subseteq
 $yRb \wedge Rx$
 $\forall b \in B, x, y \in A, B \subseteq A, R$ is a poset.

Define R on $\mathbb{Z} - \{0\}$ by aRb if $a|b$

reflexive? yes; $\forall a \in \mathbb{Z}$

antisymmetry? Does $a|b$ and $b|a$ imply $a=b$?

No: $2|1 \cdot 2$ and $-2|2$, but $2 \neq -2$.

Transitivity? Yes: If $a|b$ and $b|c$ then $a|c$, see test #1.

Def: Let (A, \prec) be a poset. Then \prec is a linear order of A if $\forall a, b \in A$, either $a \prec b$ or $b \prec a$.

Ex) (\mathbb{Z}, \leq) is a linear order.

Ex) Define \ll on \mathbb{N} by $a \ll b$ iff $a|b$

Is \ll a linear ordering?

1. Is it a poset? Yes

2. But it is not linear order: $2 \ll 3$ and $3 \ll 2$ so aR_1R_2a

Def: Let \prec be a partial order of $A \neq \emptyset$.

a) If R_1 and R_2 are irreflexive,

Prove $R_1 \cap R_2$ and $R_1 \cup R_2$ are as well.

Let $a \in R_1$. Then $a \nless a \forall a \in A$.

Let $b \in R_2$. Then $b \nless b \forall b \in A$.

and $\exists a \in R_1 \cup R_2$ s.t.

$a, a \notin R_1 \wedge a \notin R_2$

* $x \in A$ is maximal if $\nexists a \in A - \{x\}$ s.t. $x \prec a$.

* $x \in A$ is a maximum if $a \prec x \forall a \in A$

* $x \in A$ is minimal if $\nexists a \in A - \{x\}$ s.t. $a \prec x$

* $x \in A$ is a minimum if $\forall a \in A, y \prec a$.

\ll "divides"

Ex] $A \neq \emptyset$. Consider $(P(A), \subseteq)$.

minimal: \emptyset every subset of A contains \emptyset

minimum elements: \emptyset

maximal elements: A

maximal element: A (A contains every subset of $P(A)$).

Fact: Let \prec be a partial order of set $A \neq \emptyset$.

If $m \in A$ is a minimum, m is unique.

If $M \in A$ is a maximum, M is unique.

Proof Suppose m_1, m_2 are minimum elements of A w.r.t. \prec .

m_1 is a minimum $\Rightarrow m_1 \prec m_2$ } By anti-symmetry,
 m_2 is a minimum $\Rightarrow m_2 \prec m_1$ } $m_1 = m_2$

Definition: Let (A, \prec) be a linearly ordered set. If each non-empty subset of A contains a minimum element, then (A, \prec) is a well ordered set.

Ex] (\mathbb{N}, \leq) is well ordered set {well ordering Axiom.}

(\mathbb{R}, \leq) is not well ordered.

$(-\infty, 0]$: no min

$(0, 1)$: no min

Onto functions

f is onto if $\forall b \in B \exists a \in A$ such that $f(a) = b$

Least Upper bounds // Greatest lower bound.

4.1: Functions - Basic Definitions

Def: x only paired with y .

Prove F is a function:

Suppose $(x, y) \in F$ and $(x, y_2) \in F$

$(x, y_1) \in F \Rightarrow y_1 = 2x - 13$

$(x, y_2) \in F \Rightarrow y_2 = 2x - 13$ ■

$f: A \rightarrow B$

Function f maps A to B .

$\text{Dom}(f) = A ; \text{Ran}(f) \subseteq B$.

Cardinality of Sets:

Two sets A and B have the same cardinality if \exists bijection $f: A \rightarrow B$. A is countable if \exists injection $f: A \rightarrow \mathbb{N}$

Fact: \mathbb{Z} is countable. In fact, $|\mathbb{Z}| = |\mathbb{N}|$

Proof: Define $f: \mathbb{Z} \rightarrow \mathbb{N}$ by $f(z) = \begin{cases} -2z & \text{if } z < 0 \\ 2z+1 & \text{if } z \geq 0 \end{cases}$

Fact: \mathbb{R} is not countable

"Proof" suppose \mathbb{R} is countable. Then \exists bijection $f: \mathbb{R} \rightarrow \mathbb{N}$. Then you can "number" the real #'s. In fact, you could number the reals in $(0, 1)$.

Definition 4.15. A function $f: A \rightarrow B$ that is both one-to-one and onto is a **one-to-one correspondence** (or **bijection**).

Definition 4.16. Let A be any nonempty set and define $I_A(x) = x$ for all $x \in A$. This function is a one-to-one correspondence and is called the **identity function** on A .

We will take the opportunity here to introduce some notation that will be studied in greater detail in Chapter 5. $\mathcal{F}(A)$ will represent the set of all functions from a nonempty set A to itself. $\mathcal{S}(A)$ will be used to represent the subset of $\mathcal{F}(A)$ whose elements are one-to-one and onto. $\mathcal{S}(A)$ is always nonempty since the identity function on A is an element of the set (since $A \neq \emptyset$).

Example 4.17. Let $A = \{a, b\}$. Define all functions on A as follows:

\rightarrow	I	f	g	h
a	a	a	b	b
b	b	a	b	a

Table 7: $\mathcal{F}(A)$ where $A = \{a, b\}$

We have $\mathcal{F}(A) = \{I, f, g, h\}$ and $\mathcal{S}(A) = \{I, h\}$ (Here, the table is read that $f = \{(a, a), (b, a)\}$, etc.)

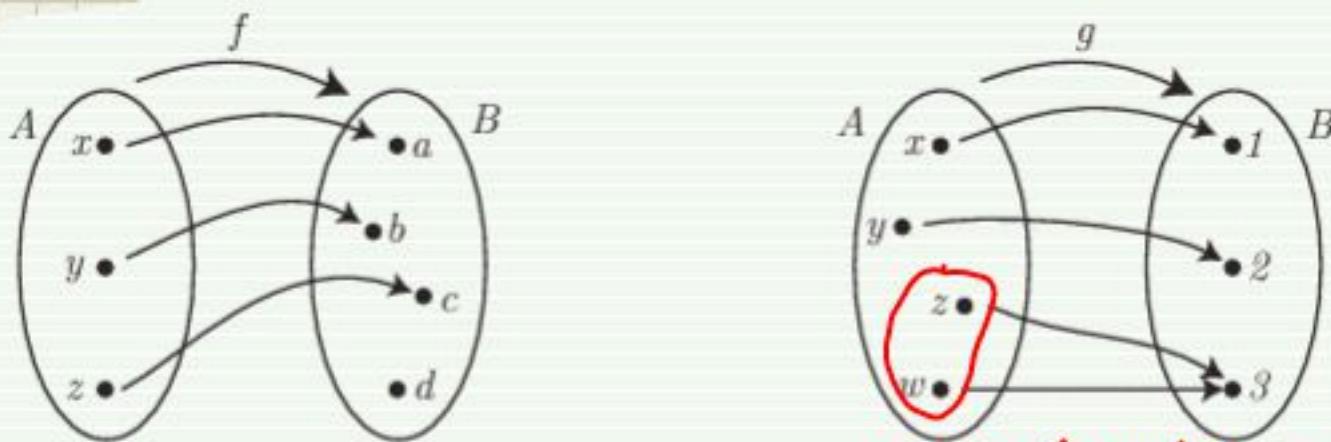


Figure 13

CAUSES problems

In Figure 13 we see that $f : A \rightarrow B$ and the arrows of f can be reversed so that f^{-1} is a function from B to A ; however, f^{-1} is not defined for all values of B and so we could not write $f^{-1} : B \rightarrow A$. If we restrict the domain, we could write $f^{-1} : \{a, b, c\} \rightarrow A$ or $f^{-1} : \text{Ran}(f) \rightarrow A$.

With the function g indicated in 13 we encounter a more serious problem. Notice that z and w are both mapped to 3 by the function g . Therefore, reversing the direction of the arrows would map 3 to both z and w , a violation of the definition of a function. **Thus it is clear that while all functions have inverses, the inverses may not be functions.** We will consider the properties a function must have to possess an inverse function. We will call such a function **invertible**.

Global Hypothesis

Theorem 4.18. Let $f : A \rightarrow B$. Then $f^{-1} : \text{Ran}(f) \rightarrow A$ if and only if f is one-to-one.

Proof. Suppose $f : A \rightarrow B$ and recall that this means f is a function, $\text{Dom}(f) = A$, and $\text{Ran}(f) \subseteq B$.

(\Leftarrow): Assume f is one-to-one. To prove that $f^{-1} : \text{Ran}(f) \rightarrow A$, we must establish that f^{-1} is a function, $\text{Dom}(f^{-1}) = \text{Ran}(f)$, and $\text{Ran}(f^{-1}) \subseteq A$.

- First, to show that f^{-1} is a function, suppose (b, a_1) and (b, a_2) are elements of f^{-1} . Then (a_1, b) and (a_2, b) are elements of f . Since f is one-to-one, $a_1 = a_2$. Therefore f^{-1} is a function.
- Note that since f is a relation we have $\text{Dom}(f^{-1}) = \text{Ran}(f)$ (by Fact 3.10).
- Finally, since f is a relation, we have $\text{Ran}(f^{-1}) = \text{Dom}(f) = A$, so $\text{Ran}(f^{-1}) \subseteq A$ (also by Fact 3.10).

(\Rightarrow): (we used double containment) \rightsquigarrow Tegrity

(\Rightarrow): Assume $f^{-1} : \text{Ran}(f) \rightarrow A$. Then f^{-1} is a function. To prove that f is one-to-one, suppose $(a_1, b), (a_2, b) \in f$. Then $(b, a_1), (b, a_2) \in f^{-1}$, and by virtue of f^{-1} being a function we can conclude that $a_1 = a_2$. Therefore, f is one-to-one. \square

f^{-1} is a function: Suppose $b_1, b_2 \in \text{Dom}(f^{-1})$ with $b_1 = b_2$

$b_1 \in \text{Dom}(f^{-1}) \rightarrow \exists a_1 \in \text{Ran}(f^{-1}) \subseteq A \text{ s.t. } (b_1, a_1) \in f^{-1}$

$b_2 \in \text{Dom}(f^{-1}) \rightarrow \exists a_2 \in \text{Ran}(f^{-1}) \subseteq A \text{ s.t. } (b_2, a_2) \in f^{-1}$

$(a_1, b_1) \in f$ and $(a_2, b_2) \in f$. $b_1 = b_2$ and f is 1-1. $\rightarrow a_1 = a_2 \rightarrow f(b_1) = f(b_2)$

4.3 Combining functions, relations

Let R be a relation from A to B and let S be a relation from B to C .

The composition of R and S , $R \circ S$, is the relation from A to C defined by

$$S \circ R = \{(a, c) \in A \times C \mid \exists b \in B \text{ with } (a, b) \in R \text{ and } (b, c) \in S\}$$

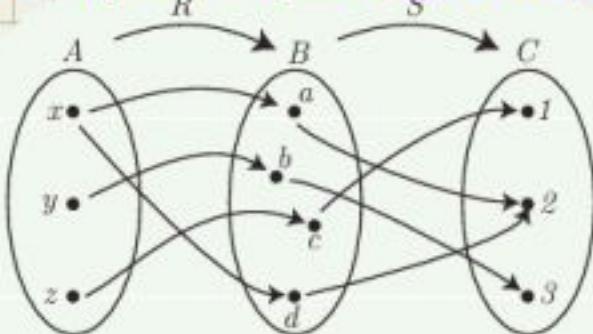


Figure 14: R is a relation from A to B and S is a relation from B to C .

Fact: Suppose $f: A \rightarrow B$ and $g: B \rightarrow C$.
Then $g \circ f: A \rightarrow C$

"Proof": Let $f: A \rightarrow B$ and $g: B \rightarrow C$. To prove $g \circ f: A \rightarrow C$, we must show:

- (1) $\text{Dom}(g \circ f) = A$,
- (2) $\text{Ran}(g \circ f) \subseteq C$
- (3) $g \circ f$ is a function.

Definition 4.20. Let R be a relation from A to B and S be a relation from B to C . Then the **composition** of R and S , denoted $S \circ R$, is defined by

$$S \circ R = \{(a, c) \in A \times C \mid \exists b \in B \text{ with } (a, b) \in R \text{ and } (b, c) \in S\}.$$

Theorem 4.27. Let $f: A \rightarrow B$ and $g: B \rightarrow C$. Then

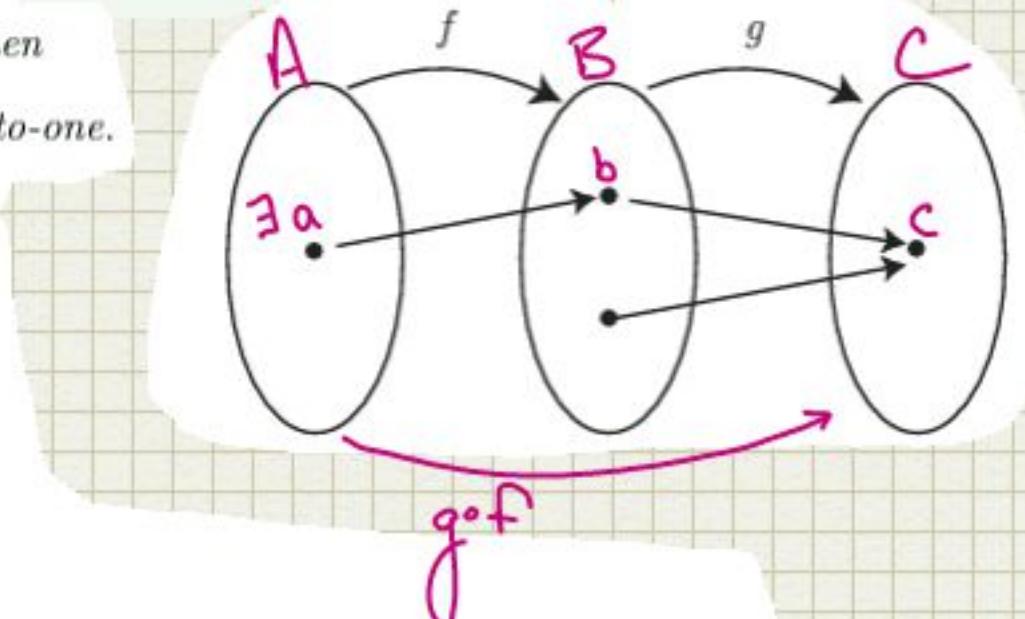
1. If f and g are one-to-one, then $g \circ f$ is one-to-one.
2. If f and g are onto, then $g \circ f$ is onto.

Theorem 4.29. Let $f: A \rightarrow B$ and $g: B \rightarrow C$.

1. If $g \circ f$ is one-to-one, then f is one-to-one.
2. If $g \circ f$ is onto, then g is onto.

Proof.

1. Assume $g \circ f$ is one-to-one. To see f is one-to-one, let $f(a_1) = f(a_2)$. Then since g is a function, we have $g(f(a_1)) = g(f(a_2))$. This implies $(g \circ f)(a_1) = (g \circ f)(a_2)$. Since $g \circ f$ is one-to-one, we have $a_1 = a_2$. Therefore, f is one-to-one.
2. Assume $g \circ f$ is onto. To see g is onto, let $c \in C$. Since $g \circ f$ is onto, there exists some $a \in A$ such that $(g \circ f)(a) = c$, and $(g \circ f)(a) = g(f(a))$. But $f(a) = b$ for some $b \in B$ and for that b we have $g(b) = c$. Therefore, g is onto



Recall $\mathcal{S}(A) = \{f: A \rightarrow A \mid F \text{ is bijective}\}$

some nice facts about \mathcal{S}

1. \circ is a binary operation

2. \circ is associative

3. \circ has a special identity

7. Let $f : \mathbb{R} \rightarrow \mathbb{R}$. We say f is strictly increasing if $a < b$ implies $f(a) < f(b)$. Prove that if f is strictly increasing, then f is one-to-one and f^{-1} is strictly increasing.

Suppose f is increasing. Let $x_1, x_2 \in \text{Dom}(f) = \mathbb{R}$ s.t. $f(x_1) = f(x_2)$.

Suppose $x_1 \neq x_2$, and WLOG suppose $x_1 < x_2$. f increasing implies $f(x_1) < f(x_2)$ #

If $y_1 < y_2$, then $f^{-1}(y_1) < f^{-1}(y_2)$

$(y_1, x_1), (y_2, x_2) \in f^{-1} \rightarrow (x_1, y_1), (x_2, y_2) \in f$.

We have $y_1 < y_2$ and f is inc.
if $x_1 \geq x_2 \xrightarrow{f \text{ inc}} y_2 \geq y_1$ #
so $x_1 < x_2$.

13) $f: A \rightarrow B$ $g: B \rightarrow C$ $g \circ f$ is onto and g is 1-1 $\rightarrow f$ is onto

Proof: Let $b \in B = \text{Dom } g \Rightarrow \exists c \in C$ s.t. $(b, c) \in g$.

$c \in C$ and $g \circ f$ onto $\Rightarrow \exists a \in A$ s.t. $(a, c) \in g \circ f$.

$(a, c) \in g \circ f \Rightarrow \exists \hat{b} \in B$ s.t. $(a, \hat{b}) \in f$ and $(\hat{b}, c) \in g$.

$(b, c), (\hat{b}, c) \in g$ and g 1-1 $\Rightarrow \underline{\underline{b = \hat{b}}}$. So $(a, b) \in f$

Thus f is onto.

Fact 4.30. Let $f: A \rightarrow B$ and $g: B \rightarrow C$. If $g \circ f$ is 1-1 and f is onto, then g must be 1-1.

Fact 4.31. Let $f: A \rightarrow B$ and $g: B \rightarrow C$. If $g \circ f$ is onto and g is 1-1, then f must be onto.

4.4 More Properties of functions/relations

Fact 4.35: Let R be a relation from A to B and S a relation from B to C . Then $(S \circ R)^{-1} = R^{-1} \circ S^{-1}$.

Proof: {double containment}

\subseteq Let $(c, a) \in (S \circ R)^{-1}$. Then $(a, c) \in S \circ R$.

By def. of \circ , $\exists b \in B$ s.t. $(a, b) \in R$ and $(b, c) \in S$.

Thus, $(b, a) \in R^{-1}$ and $(c, b) \in S^{-1}$. Since $(c, b) \in S^{-1}$ and $(b, a) \in R^{-1}$, then
by def. of \circ , $(c, a) \in R^{-1} \circ S^{-1}$.

\supseteq Let $(c, a) \in R^{-1} \circ S^{-1}$

Fact 4.36: Let R be a relation on R . R is symmetric iff $R = R^{-1}$

② R is transitive iff $R \circ R \subseteq R$.

Proof: ②

\Rightarrow Suppose R is transitive and let $(a_1, a_3) \in R \circ R$. So $\exists a_2 \in A$ s.t. $(a_1, a_2) \in R$ and $(a_2, a_3) \in R$. By transitivity, $(a_1, a_3) \in R$.

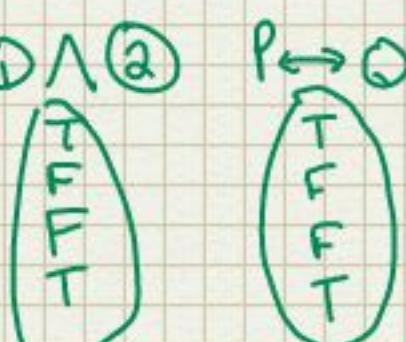
Therefore, $R \circ R \subseteq R$

\Leftarrow suppose $R \circ R \subseteq R$. Further suppose (a, b) and $(b, c) \in R$. By def. of \circ , $(a, c) \in R \circ R \subseteq R$. So $(a, c) \in R$. R is transitive.

P	Q	$(P \xrightarrow{①} Q)$	$(Q \xrightarrow{②} P)$
T	T	T	T
T	F	F	T
F	T	T	F
F	F	T	T

$$\textcircled{1} \cap \textcircled{2} \leftrightarrow (P \leftrightarrow Q)$$

P	Q	R	$\textcircled{1} \wedge \textcircled{2}$	$P \leftrightarrow Q$	T
T	T	T	T	T	T
T	T	F	F	F	T
T	F	F	F	F	T
F	T	T	F	F	T
F	F	F	F	F	T



Let $f: A \rightarrow B$, $T_1, T_2 \subseteq B$

$$\rightarrow f^{-1}(T_1 \cap T_2) = f^{-1}(T_1) \cap f^{-1}(T_2).$$

Proof: Let $a \in f^{-1}(T_1 \cap T_2)$. Then $\exists t \in T_1 \cap T_2$ s.t. $(a, t) \in f$.

$$(\subseteq) t \in T_1 \cap T_2 \rightarrow t \in T_1 \text{ and } t \in T_2. (a, t) \in f \text{ and } t \in T_1 \rightarrow a \in f^{-1}(T_1), a \in f^{-1}(T_2)$$
$$\rightarrow a \in f^{-1}(T_1) \cap f^{-1}(T_2).$$

(\supseteq) suppose $a \in f^{-1}(T_1) \cap f^{-1}(T_2)$.

$$\rightarrow a \in f^{-1}(T_1) \text{ and } a \in f^{-1}(T_2).$$

$$a \in f^{-1}(T_1) \rightarrow \exists t_1 \in T_1 \text{ s.t. } (a, t_1) \in f$$

$$a \in f^{-1}(T_2) \rightarrow \exists t_2 \in T_2 \text{ s.t. } (a, t_2) \in f$$

$(a, t_1), (a, t_2) \in f$ and f is a function $\rightarrow t_1 = t_2$. Let $t = t_1 = t_2$.

Now $(a, t) \in f$ and $t \in T_1 \cap T_2$. Since $t = t_1 \in T_1$, $t = t_2 \in T_2$. $a \in f^{-1}(T_1 \cap T_2)$

4.37 (10): $(f \circ f^{-1})(T_1) \subseteq T_1$ with $f: A \rightarrow B$, $T_1 \subseteq B$.

Proof: Let $b \in (f \circ f^{-1})(T_1)$. Then $\exists t \in T_1$ s.t. $(t, b) \in f \circ f^{-1}$. By def. of \circ . $\exists a \in A$ s.t. $(t, a) \in f^{-1}$ and $(a, b) \in f$. Then $(a, t) \in f$. So $t = b$. Since $t \in T_1$, so is b . ■

Fact 4.39 stuff.

Chapter 5 - Binary Operations.

5.1 Basic Definitions

Def 5.1: Let S be a non-empty set. If $*: S \times S \rightarrow S$, then $*$ is called a Binary Operation on S .

Ex] "+" is a binary op. on \mathbb{Z}

Notation: If $(S_1, S_2) \in S \times S$, $*((S_1, S_2)) = S_1 * S_2$

Define $*$ on \mathbb{Q} by $\frac{a}{b} * \frac{c}{d} = \frac{a+c}{b}$. Is it a binary operation on \mathbb{Q} ?

No: $*$ is not a function: $\frac{1}{2} * \frac{1}{1} = \frac{1+1}{2} = 1$ but $\frac{2}{4} * \frac{3}{3} = \frac{3+2}{4} = \frac{5}{4} \neq 1$

Fact 5.5. Each of the following is true.

1. $+$ is a binary operation on \mathbb{Z} ; that is, $+: \mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z}$.
2. \cdot is a binary operation on \mathbb{Z} ; that is, $\cdot: \mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z}$.
3. $-$ is a binary operation on \mathbb{Z} ; that is, $-: \mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z}$.
4. $+, -,$ and \cdot are binary operations on \mathbb{Q} .
5. $+, -,$ and \cdot are binary operations on \mathbb{R} .
6. \div is a binary operation on $\mathbb{Q} \setminus \{0\}$.
7. \div is a binary operation on $\mathbb{R} \setminus \{0\}$.
8. \cup is a binary operation on $\mathcal{P}(S)$.
9. \cap is a binary operation on $\mathcal{P}(S)$.
10. \setminus is a binary operation on $\mathcal{P}(S)$.

Before investigating properties of binary operations, one more example is needed. As we will see, this example has some very interesting properties and will be extremely important in later mathematics courses.

Definition 5.6. If A is a nonempty set,

$$\mathcal{S}(A) = \{f: A \rightarrow A \mid f \text{ is a bijection}\}$$

is called the **symmetric group** on A . The elements of $\mathcal{S}(A)$ are called **permutations** of A . In the special case where $A = \{1, 2, \dots, n\}$, $\mathcal{S}(A)$ is denoted S_n .

Fact 5.7. \circ is a binary operation on $\mathcal{F}(A)$ and on $\mathcal{S}(A)$.

Proof. Surely any two functions from A to A can be composed and will yield a function from A to A , so the range and domain of \circ are appropriate. Let $f, g, h, k \in \mathcal{F}(A)$ with $f = h$ and $g = k$, which means $f(x) = h(x)$ and $g(x) = k(x)$ for all $x \in A$. Surely then, for any $x \in A$, we have $(g \circ f)(x) = g(f(x)) = g(h(x)) = k(h(x)) = (k \circ h)(x)$. Hence, equals composed with equals are equal and we can conclude \circ is a binary operation on $\mathcal{F}(A)$.

Since Theorem 4.29 assures us that the composition of bijections from A to A is a bijection from A to A , we can conclude \circ is a binary operation on $\mathcal{S}(A)$. \square

} in class.

Defining Permutations in S_n .

Ex] To define a function $f: \{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}$, you can just specify what 1 maps to, what 2 maps to, etc. Composing Permutations

f	Double line notation	Consider
$1 \rightarrow 3$	$f = (1 \ 2 \ 3 \ 4)$	$f = (1 \ 2 \ 3 \ 4 \ 5 \ 6) \in S_6$ and $g = (1 \ 4 \ 3 \ 2 \ 6 \ 5) \in S_6$
$2 \rightarrow 2$		$f \circ g = (1 \ 2 \ 3 \ 4 \ 5 \ 6) \in S_6$
$3 \rightarrow 1$		$g \circ f = (1 \ 2 \ 3 \ 4 \ 5 \ 6) \in S_6$
$4 \rightarrow 1$	$f \in S_4$ in general if $f \in S_n$. write $f = (1, 2, \dots, n)$	<u>Inverse of permutations:</u> Just swap rows! $f: (1 \ 2 \ 3 \ 4 \ 5) \in S_5 \rightarrow f^{-1} = (5 \ 4 \ 3 \ 2 \ 1)$

Check: Using $f = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 5 & 4 & 2 & 1 & 3 \end{pmatrix} \rightarrow f^{-1} = \begin{pmatrix} 5 & 4 & 2 & 1 & 3 \\ 1 & 2 & 3 & 4 & 5 \end{pmatrix} = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 4 & 3 & 5 & 2 & 1 \end{pmatrix}$ Now $f \circ f^{-1} = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 1 & 2 & 3 & 4 & 5 \end{pmatrix} = I_5$

Cycles: $f \in S_n$ is a cycle if it acts like this: for some $a \in \{1, 2, \dots, n\}$,

$$a \xrightarrow{f} f(a) \xrightarrow{f} f^2(a) \xrightarrow{f} f^3(a) \xrightarrow{f} \dots \xrightarrow{f} f^{(k-1)}(a) \xrightarrow{f} a$$

and "fixes" all elements not in the chain.

Ex Consider $f = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ 4 & 3 & 6 & 5 & 7 & 2 \end{pmatrix} \in S_7$

Notation: $f = (2, 4, 6, 7) \in S_7$

Let $g = (4, 1, 3, 8, 6) \in S_9 \rightarrow g(4) = 1, g(1) = 3, g(3) = 8$

Def: 2 cycles $f = (a_1, a_2, a_3, \dots, a_k)$ and $g = (b_1, b_2, \dots, b_m) \in S_n$ are disjoint if $A \cap B = \emptyset$

Theorem 5.14. If $f = (a_1, a_2, \dots, a_k)$ and $g = (b_1, b_2, \dots, b_m)$ are disjoint cycles in S_n , then $f \circ g = g \circ f$. That is, disjoint cycles commute.

Theorem 5.15. Every permutation in S_n can be expressed as a product of disjoint cycles.

Ex
let

$$f = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 & 13 & 14 \\ 1 & 5 & 4 & 3 & 8 & 7 & 6 & 2 & 9 & 14 & 12 & 13 & 10 & 11 \end{pmatrix} \in S_{14}.$$

$$\rightarrow (1)(2, 5, 8), (3, 4), (6, 7), (9), (10, 14, 11, 13)$$

$$(3, 2) \circ (4, 2, 1, 9, 7) \circ (6, 5, 4) \circ (9, 5, 6, 7) \circ (1, 3, 5, 9, 4) = (1, 2) \circ (3, 5) \circ (4, 9, 6).$$

Chapter 6

Def: Let $x, y, z \in \mathbb{Z}, x \neq 0$. If $\exists n \in \mathbb{Z}$ with $y = x \cdot n$, we say x divides y or y is a multiple of x .

Notation: $x | y$.

If $x | y$ and $x | z$, we say x is a common divisor of y and z .

If $x | z$ and $y | z$, we say z is a common multiple of x and y .

Let $f: A \rightarrow B$ with $T_1, T_2 \subseteq B$

$$f^{-1}(T_1 \setminus T_2) = f^{-1}(T_1) \setminus f^{-1}(T_2).$$

Proof:

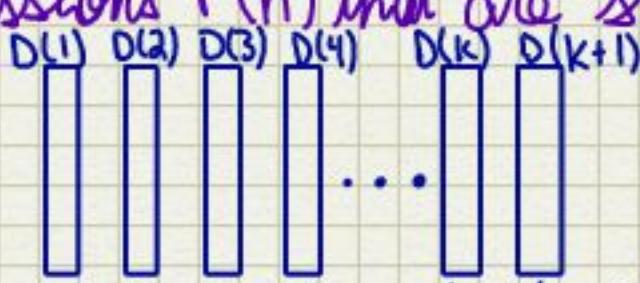
(\subseteq) Let $a \in f^{-1}(T_1 \setminus T_2) \Rightarrow \exists t \in (T_1 \setminus T_2)$ s.t. $(a, t) \in f$.

$t \in T_1 \setminus T_2 \rightarrow t \in T_1 \wedge t \notin T_2$.

Chapter 6.2: Mathematical Induction

Def: Propositional functions are expressions $P(n)$ that are statements for particular values of n .

Ex] $P(n) = "n^2 > 8"$ → $P(1) = \text{F}$ $P(2) = \text{F}$ $P(3) = \text{T}$



Ex] Let $P(n)$ be the sentence } suppose these dominoes have the
 $"1+2+3+\dots+n = \frac{n(n+1)}{2}"$ } want to prove property that for $k \in \mathbb{N}$ in $D(k)$ falls,
Axiom: (Well Ordering Principle): (\mathbb{N}, \leq) is a well-ordered set. Every subset } $D(k+1)$ falls down. } $D(1)$ falls down.

$B \subseteq \mathbb{N}$ has a minimum element.

First Principle of Mathematical Induction

Let $P(n)$ be a statement for each $n \in \mathbb{Z}$ with $n \geq 1$. If:

1. $P(1)$ is true and base case

2. $P(k)$ being true implies $P(k+1)$ is true for $k \geq 1$ $P(k) \rightarrow P(k+1)$

then $P(n)$ is true for all $n \in \mathbb{Z}$ with $n \geq 1$. Inductive Step

assume-inductive hypothesis

The Second Principle of Mathematical Induction

AKA "Strong Induction"

Let $P(n)$ be a propositional function with $n \in \mathbb{Z}$, $n \geq 1$.

If ① $P(1)$ is true, and Inductive Hypotheses

② The truth of $P(1), P(2), P(3), \dots$; and $P(k-1)$ implies the truth of $P(k)$,

Then $P(n)$ is true for all $n \in \mathbb{Z}$, $n \geq 1$.

Ex] $24 | (2 \cdot 7^n + 3 \cdot 5^n - 5) \forall n \in \mathbb{N}$.

Proof by induction:

Base case: $n=1 \rightarrow 14 + 15 - 5 = 24$; $24 | 24$ ✓ Evidently true

Inductive step: Suppose $24 | (2 \cdot 7^m + 3 \cdot 5^m - 5)$ for $m \in \mathbb{Z}$, $1 \leq m \leq k$.

Consider $(2 \cdot 7^{k+1} + 3 \cdot 5^{k+1} - 5)$

6.2 Exercises

$$(e) 1^2 - 2^2 + 3^2 - 4^2 \dots = (-1)^{n-1} \frac{n(n+1)}{2}$$

Base case

Let $n=1$

$$\begin{cases} (-1)^{(1-1)} = -1^0 = 1 = 1^2 \checkmark \\ = \frac{1(2)}{2} = 1 \end{cases} \text{ W.T.S.}$$

Inductive step: (True for k) \rightarrow (True for $k+1$)

$$\text{Suppose } 1^2 - 2^2 + 3^2 - \dots + n^2 = (-1)^{n-1} \frac{n(n+1)}{2}$$

$$\begin{aligned} \text{Then } 1^2 - 2^2 + 3^2 - \dots + (-1)^{n-1} n^2 + (-1)^{n+1} (n+1)^2 &= (-1)^{n-1} \left(\frac{n(n+1)}{2} \right) + (-1)^{n+1} (n+1)^2 \\ &= -1^{n-1} \left(\frac{n^2+n}{2} \right) + -1^n (n^2+2n+1) \end{aligned}$$

$$\frac{(-1)^{(n+1)-1}(n+1)(n+1+1)}{2}$$

$$= \frac{-1^n(n+1)(n+2)}{2} \quad (-1)^{k-1} = (-1)^k (-1)^{-1} \\ = -(-1)^k$$

$$\begin{aligned} &= -1^n \left[(-1^{-1}) \left(\frac{n^2+n}{2} \right) + (n^2+2n+1) \right] \\ &= (-1^n)(-1^{-1}) \end{aligned}$$

Prove $3 \mid (n^3 + 2n)$

Base case \rightarrow Let $n=1$ Inductive

$3 \mid (1^3 + 2(1))$ $\xrightarrow{\text{True for } k \rightarrow \text{True for } (k+1)}$

$3 \mid (1+2) \checkmark$ Let $3 \mid (k^3 + 2k)$

Then $\exists \lambda \in \mathbb{Z}$ s.t. $(k^3 + 2k) = 3\lambda$

$$k+1(k^2 + 2k+1) \quad 3 \mid ((k+1)^3 + 2(k+1))$$

$$\begin{aligned} k^3 + 2k^2 + k &\quad \xrightarrow{+} 3 \mid [k^3 + 3k^2 + 3k + 1] + (2k+2) \\ + k^2 + 2k + 1 & \\ k^3 + 3k^2 + 3k + 1 & \quad 3 \mid k^3 + 3k^2 + 5k + 3 \end{aligned}$$

Prove: For all integers $n \geq 1$, $1^2 + 2^2 + 3^2 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6}$

Proof (by induction)

base case: ($n=1$)

If $n=1$, $1^2 = \frac{1(1+1)(2 \cdot 1 + 1)}{6}$ ✓

Inductive step: True for $k \rightarrow$ True for $(k+1)$

Let $k \geq 1 \in \mathbb{Z}$ and suppose that

$$1^2 + 2^2 + 3^2 + \dots + k^2 = \frac{k(k+1)(2k+1)}{6} \quad \text{(W.T.S.)}$$

Now $1^2 + 2^2 + \dots + (k+1)^2 = 1^2 + \dots + k^2 + (k+1)^2$

$$= \frac{k(k+1)(2k+1)}{6} + (k+1)^2$$

$$= \frac{(k+1)[k(2k+1) + 6(k+1)]}{6}$$

$$= \frac{(k+1)[2k^2 + k + 6k + 6]}{6}$$

$$= \frac{(k+1)[2k^2 + 7k + 6]}{6}$$

$$= \frac{(k+1)(k+2)(2k+3)}{6} = \frac{(k+1)[(k+1)+1][2(k+1)+1]}{6} \quad f_k = \underbrace{m}_{\text{bases } \leq k-2 < k-1 < k}$$

So by FPMI, $1^2 + 2^2 + 3^2 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6} \quad \forall n \in \mathbb{N}$ ■

$\sup P(m)$ is true

for base $\leq m \leq k$ (WTS) $P(k)$

Comprehensive

Final: induction Set Equality
binarily open.

3 or 4
Proofs properties of relations
- sym, trans, reflexive

Prove n^2 even $\rightarrow n$ even

Suppose n is odd $\rightarrow n^2$ odd

Assume n^2 is even and n is odd. Then $(2k+1)$

$4k^2 + 4k + 1 \rightarrow 2(2k^2 + 2k) + 1$ odd.

short answer image of union = union of images

"Compound" proofs $(P \rightarrow q) \leftrightarrow (\neg q \rightarrow \neg P)$

more second half than first

$$\begin{array}{c} b \models a^2 \rightarrow b \models a \\ b \models a \rightarrow b \models a^2 \end{array}$$